PHIL 500

Arguments that SL is Not Correct

A Counterexample to Modus Ponens?

A Counterexample to Modus Tollens?

A Counterexample to Disjunctive Syllogism?

The Sorites Paradox

Why SL is Not Complete

An argument is VALID if and only if it is impossible for its premises to all be true while its conclusion is false.

• We decided to theorize about argument forms involving the following English constructions:

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- How well does SL do at its job?

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- Two properties we might want our theory to have:
  - CORRECTNESS: If SL tells us an argument is valid, then it is valid.
  - COMPLETENESS: If an argument's valid, then SL tells us it is valid

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- Because SL is not complete, we will need to look at additional kinds of logical forms. This will be the task of *predicate logic*, PL.
- First: let's consider why some think that SL is not even *correct*.

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- A word of warning: these arguments are controversial (some more so than others); and many logicians are not moved by them to reject the correctness of SL.
- There's two components to our theory SL:
- the theory about which arguments *involving the sentences* of SL are valid; and
- ▶ the translation guide from English into SL.

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  - ▷ '→' is not a perfect translation of 'if..., then...'
  - ▷ the first two arguments attempt to show that the differences between '→' and 'if..., then...' prevent the English 'if..., then...' from satisfying *modus ponens* (→ *E*) and *modus tollens*.
- The second two arguments object not to the translation guide, but rather to SL's theory about which arguments *involving the sentences of* SL are valid.

#### Arguments that SL is Not Correct

#### A Counterexample to Modus Ponens?

A Counterexample to Modus Tollens?

A Counterexample to Disjunctive Syllogism?

The Sorites Paradox

Why SL is Not Complete

If Clinton doesn't win, then, if a Democrat wins, then Bernie wins.

- Clinton doesn't win.
- : If a Democrat wins, then Bernie wins.

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Clinton doesn't win.

- : If a Democrat wins, then Bernie wins.
- McGee: this is of the form *modus ponens*

if  $\mathcal{A}$  then  $\mathfrak{B}$  $\mathcal{A}$  $\therefore \mathfrak{B}$ 

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• However, in the run up to the 2016 presidential election, its premises were true yet its conclusion was false.

• The following sentence forms are provably equivalent in SL:

 $(\mathscr{A} \land \mathscr{B}) \to \mathscr{C}$  $\mathscr{A} \to (\mathscr{B} \to \mathscr{C})$ 

If Clinton doesn't win and a Democrat wins, then Bernie wins. Clinton doesn't win.

: So, if a Democrat wins, then Bernie wins.

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- If we accept *exportation* and *modus ponens*, then we will have to say that this is also valid.

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1 
$$(\neg C \land D) \rightarrow B$$
  
2  $\neg C$   
3  $\neg C \rightarrow (D \rightarrow B)$  1, Exp  
4  $D \rightarrow B$  2, 3,  $\rightarrow E$ 

McGee shows that, if we accept both *modus ponens* and *exportation* for the English 'if..., then...', then the English 'if..., then...' will be logically indistinguishable from SL's conditional →.
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- McGee shows that, if we accept both *modus ponens* and *exportation* for the English 'if..., then...', then the English 'if..., then...' will be logically indistinguishable from SL's conditional →.
- So, we'd have to accept the following argument as deductively valid:

Shakespeare wrote Hamlet.

- ∴ If Shakespeare didn't write Hamlet, then Dan Brown did.
- McGee: this is unacceptable, so we must choose between *exportation* and *modus ponens* for the English 'if..., then...'

• Actually,

Shakespeare wrote Hamlet.

∴ If Shakespeare didn't write Hamlet, then Dan Brown did.

gives us a reason, on its own, to doubt the adequacy of our translation guide, since this is an entailment, given our translation guide.



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- Others opt for modus ponens and reject exportation.
- Still others accept both *exportation* and *modus ponens* and accept that the English 'if..., then...' is logically indistinguishable from →
  - they have stories to tell about why arguments like the ones above appear—falsely—to be invalid

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• Imagine that we have an urn which contains 100 marbles.

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• Suppose that we have selected a marble at random from the urn, but that we do not yet know whether it is blue or red, or whether it is big or small.

If the marble is big, then it's likely red.

The marble is not likely red.

 $\therefore$  The marble is not big.

If the marble is big, then it's likely red. The marble is not likely red.

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If the marble is big, then it's likely red. The marble is not likely red.

- $\therefore$  The marble is not big.
- This is an instance of the argument form *modus tollens*

If  $\mathcal{A}$  then  $\mathcal{B}$ 

It is not the case that  $\mathscr{B}$ 

 $\therefore$  It is not the case that  $\mathscr{A}$ 

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If A then <del>B</del> It is not the case that <del>B</del>

- $\therefore$  It is not the case that  $\mathscr{A}$
- So, Yalcin contends, *modus tollens* is not valid for the English 'if..., then...'

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- Some options:
  - We could contend that the proposed counterexample equivocates with respect to 'likely'
    - in the first premise, it means "likely given all the information that currently have, plus the information that the marble is big"
    - ▷ in the second premise, it means "likely, given all the information that we currently have".
  - We could contend that the first premise is equivalent to "it's likely that, if the marble is big, then it's red", so that "if..., then..." is not the main operator of the premise.

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• Consider the statement:

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This very statement is false.

• Consider the statement: L := L is false. • Consider the statement: I := I is false

L := L is false.

• Priest: *L* is both true and false. (Almost everyone else disagrees)

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  - ▶ true (and not false) 'T'
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- updated truth-table for '¬':

A	¬A
Т	F
F	Т
В	В

• Updated truth-tables for  $\lor$  and  $\rightarrow$ :



• Let *P* := Pigs can fly

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$L \lor P$	$L \rightarrow P$
$\neg L$	L
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• Priest: for each argument, the premises are all true (and also false), yet the conclusion is false.
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- Priest: for each argument, the premises are all true (and also false), yet the conclusion is false.
- Yet these arguments are of the form *disjunctive syllogism* and *modus ponens*
- So, Priest concludes: both *disjunctive syllogism* and *modus ponens* are invalid.

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L' := L' is not true.

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- Any pair of sequential tiles are perceptually indistinguishable.
- By the end of the sequence, we have a tile that is unmistakably orange.

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• Something's gone wrong with this reasoning, and some people have been tempted to point the finger at *modus ponens*.

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The Sorites Paradox

Why SL is Not Complete

Johann knows Filipa.

∴ So, somebody knows Filipa.

Everyone who owns a Ford owns a car. Rohan owns a Ford.

 $\therefore$  So, Rohan owns a car.

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... So, somebody knows Filipa.

Everyone who owns a Ford owns a car. Rohan owns a Ford.

- $\therefore$  So, Rohan owns a car.
- Both arguments are valid

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- : So, Rohan owns a car.
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- Neither arguments are entailments
  - ⊳ J∴ S
  - $\triangleright \ E \ , \ F \ . \ C$