

Sentential Logic

Syntax and Semantics

PHIL 500

$$\begin{array}{c} (\neg P \wedge Q) \\ \wedge \\ \neg P \quad Q \\ | \\ P \end{array}$$

$$\begin{array}{c} \neg(P \wedge Q) \\ | \\ (P \wedge Q) \\ \wedge \\ P \quad Q \end{array}$$

The Language SL

Syntax for SL

Semantics for SL

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The Plan

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- Provide a method for translating statements of English into SL and statements of SL into English
- The advantage: we can theorize about relations of deductive validity without having to worry about the ambiguity of English

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- a way to interpret the *meaning* of every grammatical expression of the language
 - *e.g.*, a dictionary entry for every word of English and rules for constructing the meaning of sentences out of the meanings of words

Languages in General

SYNTAX — { 1. Vocabulary
2. Grammar
SEMANTICS — 3. Meaning

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- Nothing else is included in the vocabulary of SL.

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\leftrightarrow) If ' \mathcal{A} ' and ' \mathcal{B} ' are sentences, then ' $(\mathcal{A} \leftrightarrow \mathcal{B})$ ' is a sentence.

¬) Nothing else is a sentence.

- ▶ ' \mathcal{A} ' and ' \mathcal{B} ' do not appear in the vocabulary of SL.

Meta-variables

- ▶ ‘ \mathcal{A} ’ and ‘ \mathcal{B} ’ do not appear in the vocabulary of SL.
- ▶ They are *metavariables*—variables whose potential values are sentences of SL.

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 - d) So, ' $\neg(P \vee Q)$ ' is a sentence [from (c) and (\neg)]
 - e) ' R ' is a sentence [from (SL)]
 - f) So, ' $\neg(P \vee Q) \rightarrow R$ ' is a sentence [from (d), (e), and (\rightarrow)]

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Main Operators

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- So, the main operator is ' \wedge '

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 - d) So, ' $\neg(P \wedge Q)$ ' is a sentence [from (c) and (\neg)]
- ' $\neg(P \wedge Q)$ ' is not the same sentence as ' $(\neg P \wedge Q)$ '

Subsentences

- ‘ \mathcal{B} ’ is a *subsentence* of ‘ \mathcal{A} ’ if and only if, in the course of building up ‘ \mathcal{A} ’ by applying the rules for sentences, ‘ \mathcal{B} ’ appears on a line before ‘ \mathcal{A} ’.

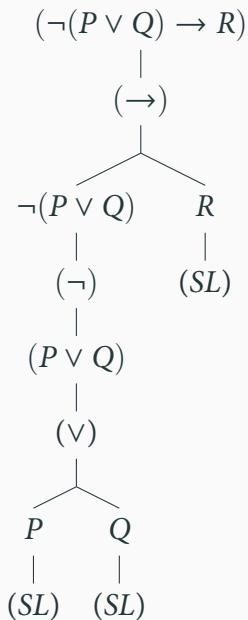
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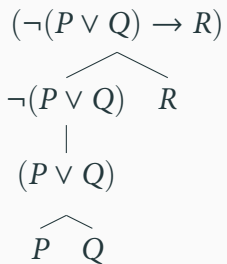
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 - ‘ $\neg P$ ’ is a subsentence of ‘ $\neg P \wedge Q$ ’
 - ‘ $\neg P$ ’ is *not* a subsentence of ‘ $\neg(P \wedge Q)$ ’

Syntactic Structure



Syntactic Structure



Syntactic Structure

$(\neg P \wedge Q)$

$\neg P \quad Q$

P

$\neg(P \wedge Q)$

$(P \wedge Q)$

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- Parentheses just make syntactic structure explicit. They do not make any contribution to meaning beyond that.
- So: we must say what the meanings of the statements letters are and what the meanings of the logical operators are

The Meaning of the Statement Letters

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- The statement letter is true if and only if the statement in English is true.

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 - we are using ' \mathcal{A} ' and ' \mathcal{B} ' as variables ranging over the sentences of SL

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F	T	T
F	F	

The Meaning of '∨'

- A sentence whose main operator is '∨' is known as a *disjunction*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \vee \mathcal{B}$
T	T	T
T	F	T
F	T	T
F	F	F

The Meaning of ' \rightarrow '

- A sentence whose main operator is \rightarrow is known as a *conditional*.

The Meaning of '→'

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\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	
T	F	
F	T	
F	F	

The Meaning of '→'

- A sentence whose main operator is \rightarrow is known as a *conditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	
F	T	
F	F	

The Meaning of '→'

- A sentence whose main operator is \rightarrow is known as a *conditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
F	T	
F	F	

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\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
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\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
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\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
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F	T	T
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- Note: this is the only binary operator which is not symmetric.

The Meaning of '→'

- A sentence whose main operator is \rightarrow is known as a *conditional*.

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T	T	T
T	F	F
F	T	T
F	F	T

- Note: this is the only binary operator which is not symmetric.
 - ' $\mathcal{A} \rightarrow \mathcal{B}$ ' does not have the same meaning as ' $\mathcal{B} \rightarrow \mathcal{A}$ '

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	
T	F	
F	T	
F	F	

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	
F	T	
F	F	

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
F	T	
F	F	

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
F	T	F
F	F	

The Meaning of ' \leftrightarrow '

- A sentence whose main operator is ' \leftrightarrow ' is known as a *biconditional*.

\mathcal{A}	\mathcal{B}	$\mathcal{A} \rightarrow \mathcal{B}$
T	T	T
T	F	F
F	T	F
F	F	T

Determining the truth-value of a sentence of SL

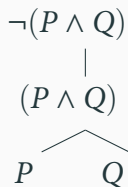
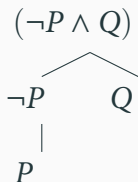
- Suppose ' P ' is true and ' Q ' is false

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- Suppose ' P ' is true and ' Q ' is false
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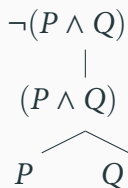
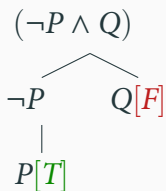
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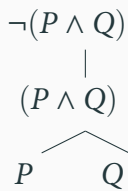
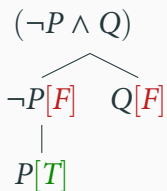
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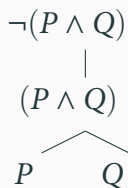
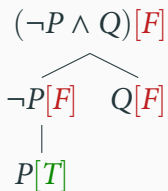
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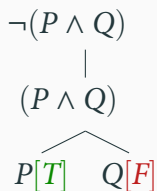
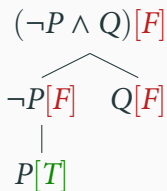
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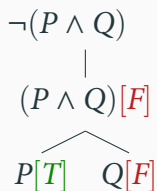
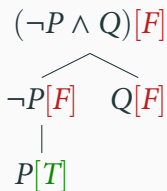
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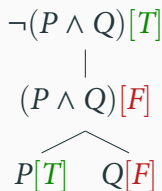
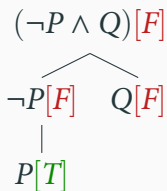
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- Suppose 'P' is true and 'Q' is false
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Determining the truth-value of a sentence of SL

- Suppose 'P' is true and 'Q' is false
- What are the truth-values of ' $\neg P \wedge Q$ ' and ' $\neg(P \wedge Q)$ '?



Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	$\neg P \wedge Q$

Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	$\neg P \wedge Q$
T	T	
T	F	
F	T	
F	F	

Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	$\neg P$	$P \wedge Q$
T	T	F	T
T	F	F	F
F	T	T	F
F	F	T	F

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- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	$\neg P$	$P \wedge Q$
T	T	F	T
T	F	F	F
F	T	T	T
F	F	T	F

Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	\neg	P	\wedge	Q
T	T	F	T	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	F	F	F

Truth Tables

- Truth-table for ' $\neg P \wedge Q$ ':

P	Q	\neg	P	\wedge	Q
T	T	F	T	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	F	F	F

Truth Tables

- Truth-table for ' $\neg(P \wedge Q)$ ':

Truth Tables

- Truth-table for ' $\neg(P \wedge Q)$ ':

P	Q	$\neg (P \wedge Q)$

Truth Tables

- Truth-table for ' $\neg(P \wedge Q)$ ':

P	Q	$\neg (P \wedge Q)$
T	T	
T	F	
F	T	
F	F	

Truth Tables

- Truth-table for ' $\neg(P \wedge Q)$ ':

P	Q	$\neg (P \wedge Q)$
T	T	T
T	F	F
F	T	F
F	F	F

Truth Tables

- Truth-table for $\neg(P \wedge Q)$:

P	Q	$\neg (P \wedge Q)$
T	T	T T T
T	F	T F F
F	T	F F T
F	F	F F F

Truth Tables

- Truth-table for $\neg(P \wedge Q)$:

P	Q	\neg	$(P$	\wedge	$Q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	T	F	F	T
F	F	T	F	F	F